

Notes on **Equitable Partitions** into **Matching Forests** in Mixed Graphs and ***b*-branchings** in Digraphs

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Equitable Coloring

[Hajnal, Szemerédi '69] (Conjecture by P. Erdős '64)

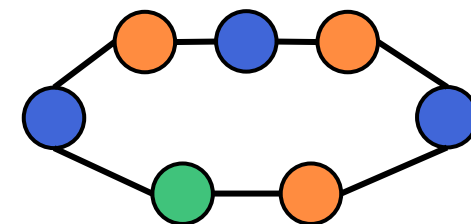
Graph (V, E) with maximum degree Δ

→ V can be partitioned into

$\Delta+1$ stable sets $S_1, S_2, \dots, S_{\Delta+1}$ such that

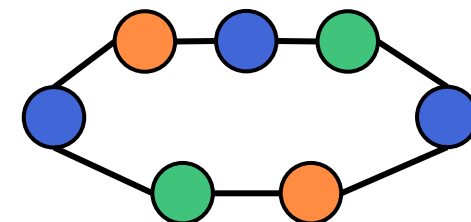
$$||S_i| - |S_j|| \leq 1 \quad \forall i, j \in \{1, 2, \dots, \Delta+1\}$$

➤ Namely, $|S_i| = \lfloor \frac{|V|}{\Delta+1} \rfloor$ or $\lceil \frac{|V|}{\Delta+1} \rceil$



(3, 3, 1)

→ Not equitable



(3, 2, 2)

→ Equitable

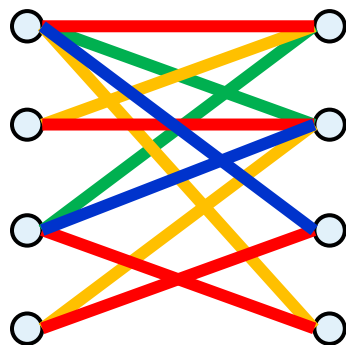
[Folkman, Fulkerson '67] etc.

Bipartite graph (V, E) with maximum degree Δ

→ For any $k \geq \Delta$, E can be partitioned into

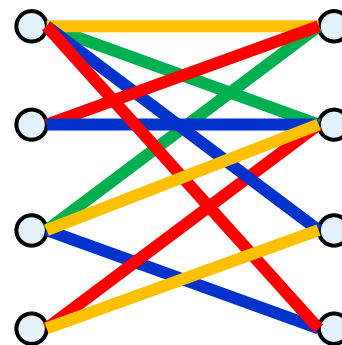
k matchings M_1, M_2, \dots, M_k such that

$$||M_i| - |M_j|| \leq 1 \quad \forall i, j \in \{1, 2, \dots, k\}$$



(4, 3, 2, 2)

→ Not equitable



(3, 3, 3, 2)

→ Equitable

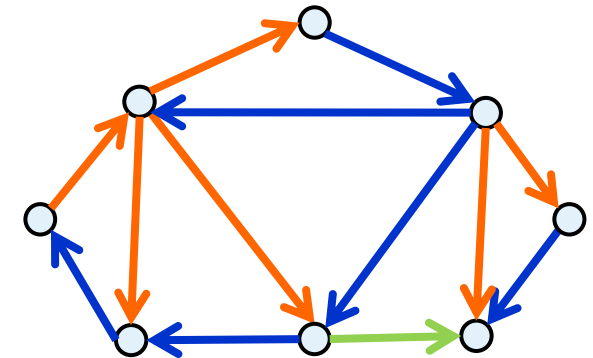
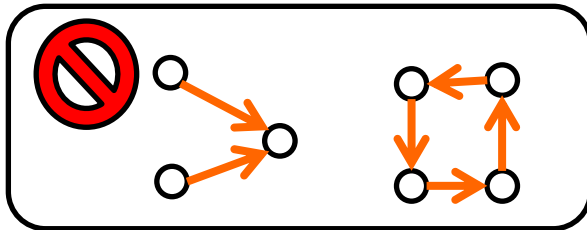
Equitable Partition into Branchings

■ Digraph (V, A)

Definition

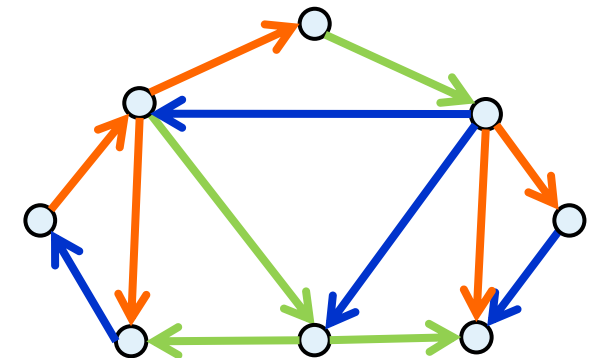
$B \subseteq A$ is a **branching** if

- (i) $\text{indeg}(v) \leq 1 \quad \forall v \in V$
- (ii) No (undirected) cycle



(6, 5, 1)

→ *Not equitable*



(4, 5, 4)

→ *Equitable*

Theorem [Schrijver 03]?

If A can be partitioned into k branchings

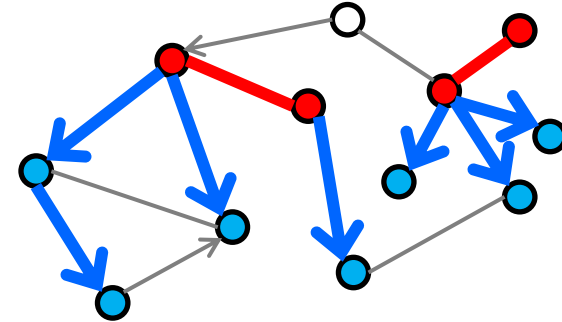
→ A can be partitioned into

k **branchings** B_1, B_2, \dots, B_k such that

$$||B_i| - |B_j|| \leq 1 \quad \forall i, j \in [k]$$

Theorem

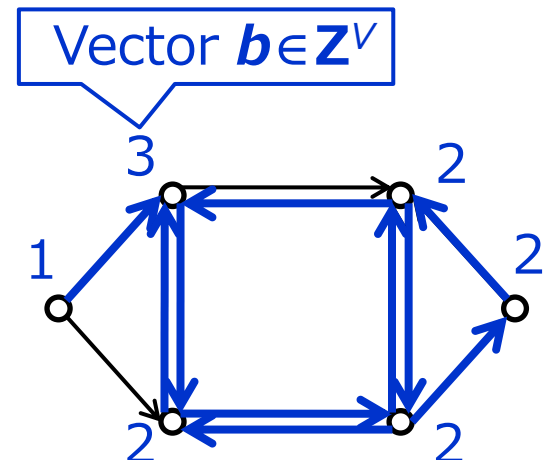
For a **mixed graph** $(V, E \cup A)$,
if $E \cup A$ can be partitioned into k **matching forests**
→ $E \cup A$ can be partitioned into
 k **matching forests** F_1, F_2, \dots, F_k such that
 $||\partial F_i| - |\partial F_j|| \leq 2 \quad \forall i, j \in \{1, 2, \dots, k\}$



Theorem

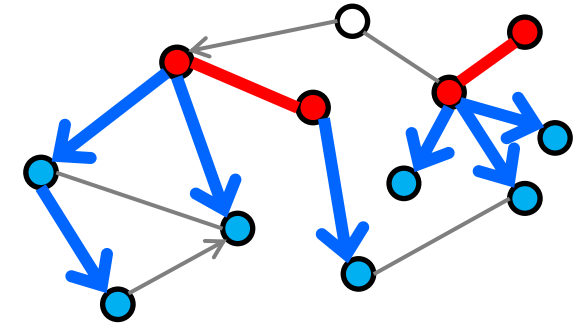
If A can be partitioned into k **b -branchings**
→ A can be partitioned into
 k **b -branchings** B_1, B_2, \dots, B_k such that

- $||B_i| - |B_j|| \leq 1 \quad \forall i, j \in [k]$
- $|\text{indeg}_i(v) - \text{indeg}_j(v)| \leq 1 \quad \forall v \in V, \forall i, j \in [k]$



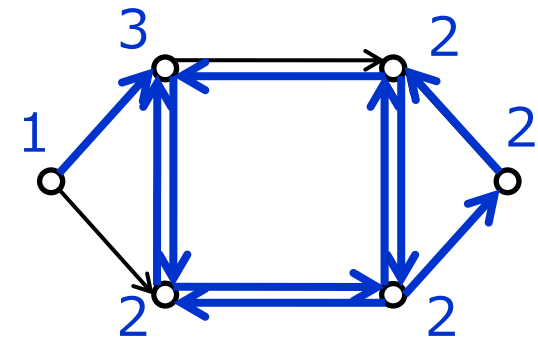
● Introduction

- Equitable partition in graphs:
 - Matching, Branching



● Matching Forest

- Common generalization of matching and branching
- [Király, Yokoi '18] **Tri-criteria equitability**
 - Cannot be optimized simultaneously
- **[This talk] Single-criterion equitability**
 - Can always be optimized



● b-branching

- Generalization of branchings allowing indegree ≥ 2
- **[This talk] $(n+1)$ -criteria equitability**
 - Can always be optimized simultaneously

● Proof Sketch (for Matching Forests)

Matching Forest [Giles '82]

- Mixed graph $G = (V, E \cup A)$

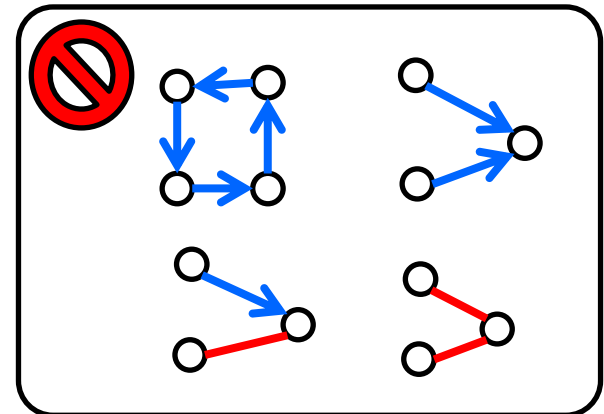
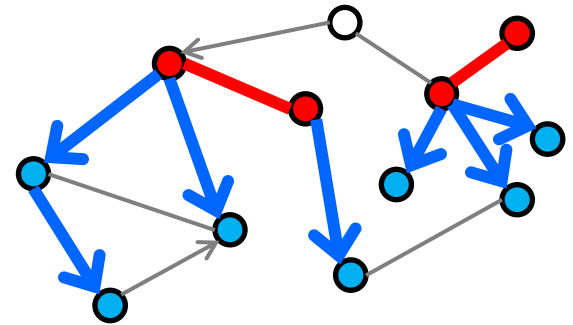
Definition

Undirected edge $\{u, v\} \in E$ **covers** both u and v
Directed edge $(u, v) \in A$ **covers** only v

Definition

$F \subseteq E \cup A$ is a **matching forest** if

- its underlying edge set is a forest
 - every vertex is covered by ≤ 1 edge
- Namely,
- $B = F \cap A$: Branching
 - $M = F \cap E$: Matching s.t. $\partial M \subseteq V \setminus \partial B$
- Set of covered vertices $\partial F := \partial B \cup \partial M$



- **Finding a max. weight matching forest** (Not a topic of today)
 - [Giles '82] Polyhedral description, Primal-dual algorithm in $O(n^2m)$ time
 - [T. '14] Improved to $O(n^3)$
 - [Schrijver '00] TDI-ness of the description
 - [Schrijver '03] Reduction to weighted linear matroid parity
- **Partition into Matching Forests**
 - [Keijsper '03] Partition into $\Delta+1$ matching forests
 - **[Király, Yokoi '18] Equitable partition into matching forests**

Theorem [Király, Yokoi '18]

If $E \cup A$ can be partitioned into k matching forests

→ $E \cup A$ can be partitioned into k matching forests F_1, F_2, \dots, F_k such that

$$||F_i| - |F_j|| \leq 1, \quad ||B_i| - |B_j|| \leq 2, \quad ||M_i| - |M_j|| \leq 2$$

→ $E \cup A$ can be partitioned into k matching forests F_1, F_2, \dots, F_k such that

$$||F_i| - |F_j|| \leq 2, \quad ||B_i| - |M_j|| \leq 2, \quad ||M_i| - |M_j|| \leq 1$$

- **Tricriteria equitability**
- These values are **tight**
i.e., the three criteria **cannot be optimized simultaneously**

- Subset family $\mathcal{F} \subseteq 2^V$

Definition

A set system (V, \mathcal{F}) is a **delta-matroid** if

$$\forall S_1, S_2 \in \mathcal{F}, \forall s \in S_1 \Delta S_2,$$

➤ $S_1 \Delta \{s\} \in \mathcal{F}$ or

➤ $\exists t \in (S_1 \Delta S_2) - \{s\}, S_1 \Delta \{s, t\} \in \mathcal{F}$

- Undirected graph $G=(V,E)$

➤ $\mathcal{F}_M := \{\partial M \subseteq V \mid M \text{ is a matching in } G\}$

Theorem [Chandrasekaran, Kabadi '88; Bouchet '89]

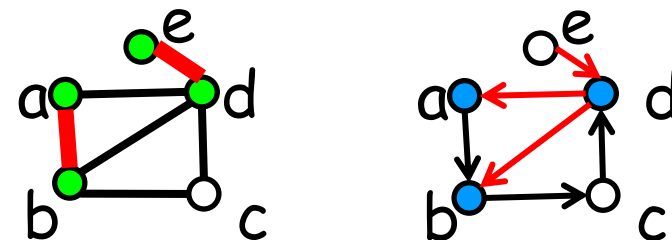
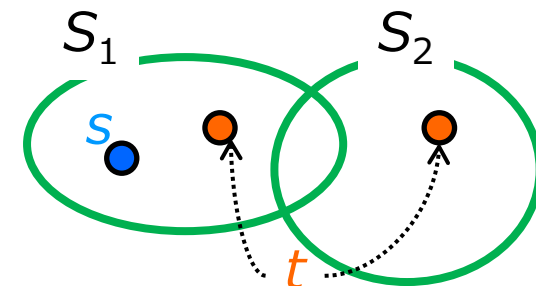
(V, \mathcal{F}_M) is a **delta-matroid**

- Directed graph $G=(V, A)$

➤ $\mathcal{F}_B := \{\partial^- B \subseteq V \mid B \text{ is a branching in } G\}$

Theorem [T. '14] (Folklore?)

(V, \mathcal{F}_B) is a matroid (thus a **delta-matroid**)



$$\mathcal{F}_M = \{\emptyset, \{a,b\}, \{a,d\}, \{b,c\}, \{b,d\}, \{c,d\}, \{d,e\}, \{a,b,c,d\}, \{a,b,d,e\}, \{b,c,d,e\}\}$$

$$\mathcal{F}_B = \{\emptyset, \{a\}, \{b\}, \{c\}, \{d\}, \{a,b\}, \{a,c\}, \dots, \{c,d\}, \{a,b,c\}, \{a,b,d\}, \{a,c,d\}, \{b,c,d\}, \{a,b,c,d\}\}$$

Equitability from the Structure

■ Mixed graph $G=(V,E\cup A)$

➤ $\mathcal{F}_{MF} := \{\partial F \subseteq V \mid F \text{ is a matching forest in } G\}$

Theorem [T. '14]

(V, \mathcal{F}_{MF}) is a **delta-matroid**

Our Idea

**Define the equitability of matching forests
by the size of ∂F**

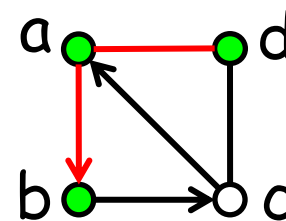
Theorem

If $E\cup A$ can be partitioned into k matching forests

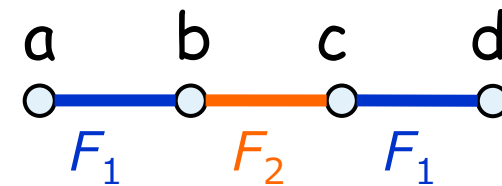
➔ $E\cup A$ can be partitioned into
 k matching forests F_1, F_2, \dots, F_k such that

$$||\partial F_i| - |\partial F_j|| \leq 2 \quad \forall i, j \in \{1, 2, \dots, k\}$$

- Value **2** is **optimal**
- Optimality always attained



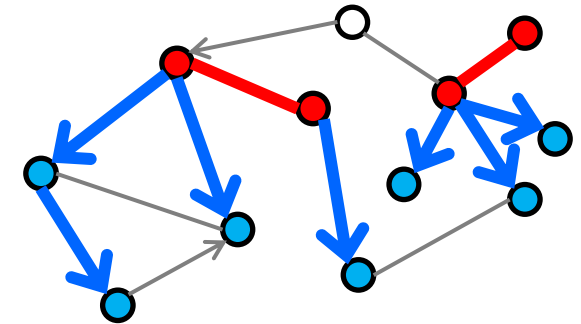
$$\mathcal{F}_{MF} = \{\emptyset, \{a\}, \{b\}, \{c\}, \\ \{a,b\}, \{a,c\}, \{a,d\}, \{b,c\}, \\ \{c,d\}, \{a,b,d\}, \{a,c,d\}, \\ \{b,c,d\}, \{a,b,c,d\}\}$$



$$\partial F_1 = \{a,b,c,d\}$$
$$\partial F_2 = \{b,c\}$$

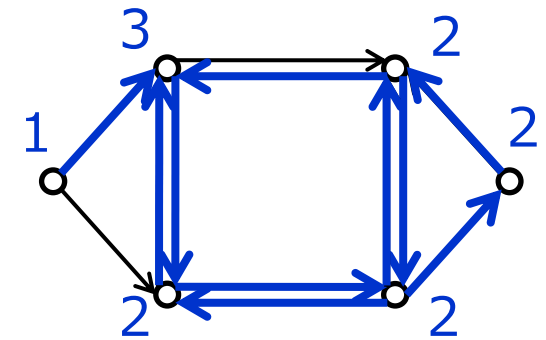
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● Matching Forest

- Common generalization of matching and branching
- [Király, Yokoi '18] **Tri-criteria equitability**
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- Generalization of branchings allowing indegree ≥ 2
- **[This talk] $(n+1)$ -criteria equitability**
 - Can always be optimized simultaneously

● Proof Sketch (for Matching Forests)

- Digraph (V, A)
- Positive integer vector $\mathbf{b} \in \mathbf{Z}^V$ on V

Definition

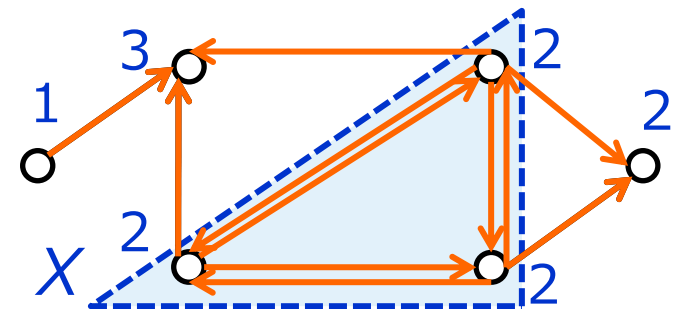
$B \subseteq A$ is a **b -branching** if

$$\begin{cases} \text{(i) } \text{indeg}_B(u) \leq b(u) & (u \in V) \\ \text{(ii) } |B[X]| \leq b(X) - 1 & (\emptyset \neq X \subseteq V) \end{cases}$$

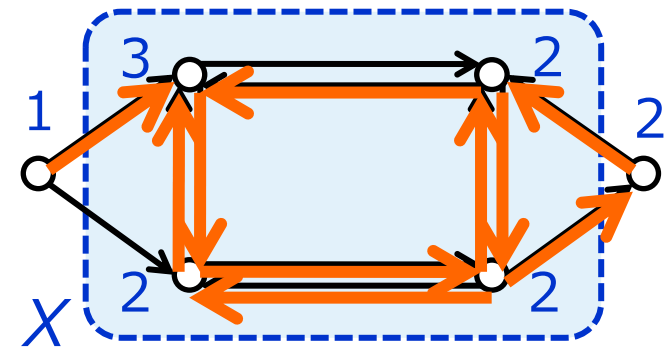
- Branching: $b(u) \equiv 1$
- Special case of **matroid intersection**
 - (i): Direct sum of uniform matroids
 - (ii): Sparsity matroid (Count matroid)

Sparsity matroid [cf. Frank's book 11]

Graph $G=(V,E)$, Vector $\mathbf{b} \in \mathbf{Z}^V$, Integer $k \geq 0$
➔ $\{B \subseteq E : |B[X]| \leq b(X) - k\}$ is an independent set family of a matroid



Not a b -branching:
 $|B[X]| = 6, b(X) - 1 = 5$



A b -branching:
 $|B[X]| = 7, b(X) - 1 = 8$

- Finding a max. weight b -branching (Not a topic of today)
 - TDI description (← Matroid intersection polytope)
 - Multi-phase greedy algorithm extending Arborescence Algorithm [Kakimura, Kamiyama, T. '20+]
- Partition into b -branchings

Theorem [Kakimura, Kamiyama, T. '20+]

A can be partitioned into k b -branchings iff

- $\text{indeg}(u) \leq k \cdot b(u) \quad \forall u \in V$
- $|A[X]| \leq k(b(X) - 1) \quad \emptyset \neq \forall X \subseteq V$

- Obvious necessary condition works

Definition (Recap)

$B \subseteq A$ is a b -branching if

- $$\left[\begin{array}{ll} \text{(i) } \text{indeg}_B(u) \leq b(u) & (u \in V) \\ \text{(ii) } |B[X]| \leq b(X) - 1 & (\emptyset \neq X \subseteq V) \end{array} \right.$$

Theorem

If A can be partitioned into k **b -branchings**

→ A can be partitioned into k

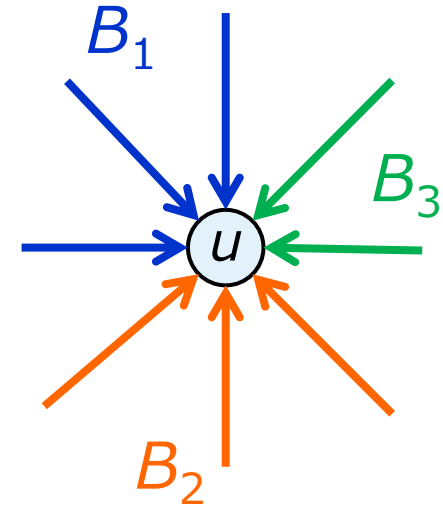
b -branchings B_1, B_2, \dots, B_k such that

- $||B_i| - |B_j|| \leq 1 \quad \forall i, j \in [k]$
- $|\text{indeg}_i(u) - \text{indeg}_j(u)| \leq 1 \quad \forall u \in V, \forall i, j \in [k]$

Namely, for each $i \in [k]$,

- $|B_i| = \lfloor \frac{|A|}{k} \rfloor$ or $\lceil \frac{|A|}{k} \rceil$
- $|\text{indeg}_i(u)| = \lfloor \frac{\text{indeg}_A(u)}{k} \rfloor$ or $\lceil \frac{\text{indeg}_A(u)}{k} \rceil \quad \forall u \in V$

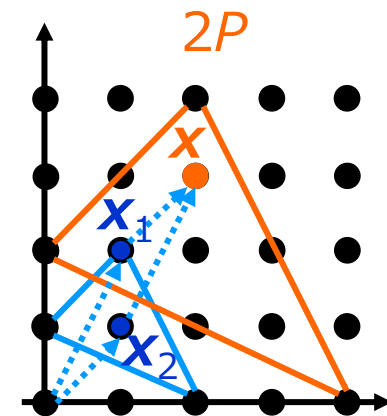
→ These **$(n+1)$ -criteria** can be simultaneously optimized



Definition

A polytope P has the **integer decomposition property** if $\forall k \in \mathbf{Z}_{++}, \forall \mathbf{x} \in kP \cap \mathbf{Z}^A, \mathbf{x} = \mathbf{x}_1 + \dots + \mathbf{x}_k$ ($\mathbf{x}_1, \dots, \mathbf{x}_k \in P \cap \mathbf{Z}^A$)

- **True** for
 - Polymatroids [Giles '75; Baum, Trotter '81]
 - Intersection of two strongly base orderable matroids [Davies, McDiarmid '76; McDiarmid '76]
 - Branching polytope (below)
- **NOT True** for every matroid intersection



Theorem [Baum, Trotter '81]

The **branching polytope** has IDP

Theorem [McDiarmid '83]

The convex hull of the incidence vectors of the **branchings of size ℓ** has IDP

Branching polytope

- $x(\delta u) \leq 1$ ($u \in V$)
- $x(A[X]) \leq |X| - 1$ ($\emptyset \neq X \subseteq V$)
- $0 \leq x(a) \leq 1$ ($a \in A$)

[Baum, Trotter '81]

Branching

[Kakimura, Kamiyama, T. '20+]

The **b -branching polytope** has **IDP**

[McDiarmid '83]

Branchings of size ℓ

[Our Result]

The convex hull of the incidence vectors of the **b -branchings of size ℓ** has **IDP**

➤ Further, the indegree can be fixed to be $b'(v)$ ($\leq b(v)$)

[Our Result]

The convex hull of the incidence vectors of

the **b -branchings** of

➤ **size ℓ** ; and

➤ **$\text{indeg}(v) = b'(v) \quad \forall v \in V$**

has **IDP**

b -branching polytope

➤ $x(\delta u) \leq b(u) \quad (u \in V)$

➤ $x(A[X]) \leq b(X) - 1 \quad (\emptyset \neq X \subseteq V)$

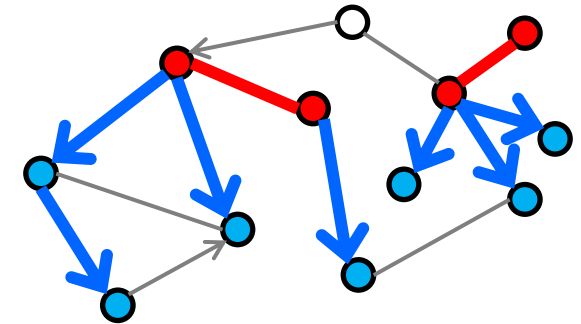
➤ $0 \leq x(a) \leq 1 \quad (a \in A)$

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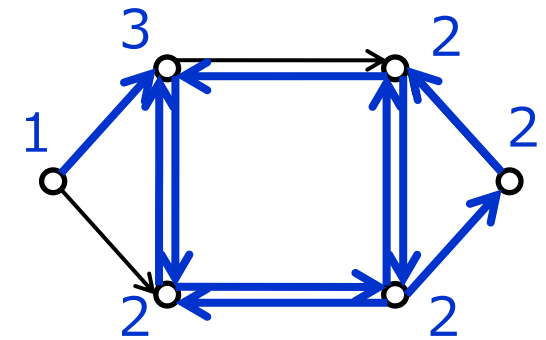
● Matching Forest

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● b-branching

- Generalization of branchings allowing indegree ≥ 2
- **[This talk] $(n+1)$ -criteria equitability**
 - Can always be optimized simultaneously



● Proof Sketch (for Matching Forests)

- Branching B
- Root set $R(B) = V \setminus \partial^- B$

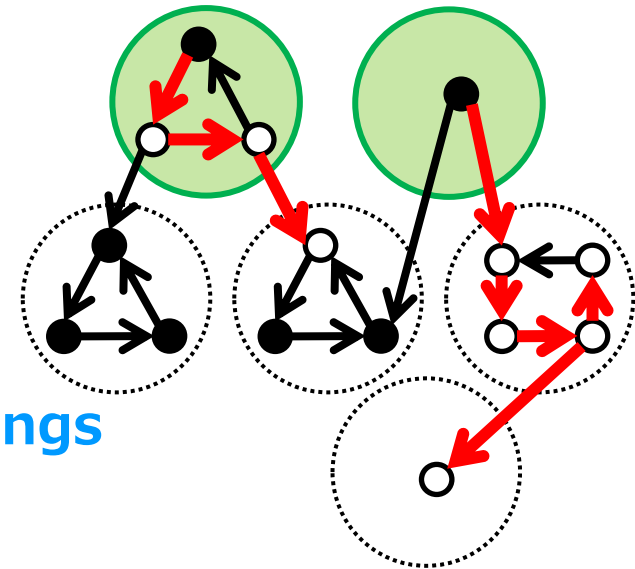
Key Lemma [Schrijver '00]

Branchings B_1, B_2 partitioning A

$X, Y \subseteq V$ satisfying $X \cup Y = R(B_1) \cup R(B_2)$, $X \cap Y = R(B_1) \cap R(B_2)$

- ➔ \exists Branchings B_1', B_2' partitioning A such that $R(B_1') = X$, $R(B_2') = Y$ iff
- $|X \cap C| \geq 1$ and $|Y \cap C| \geq 1$
 - for each **strong component** C w/o incoming arcs

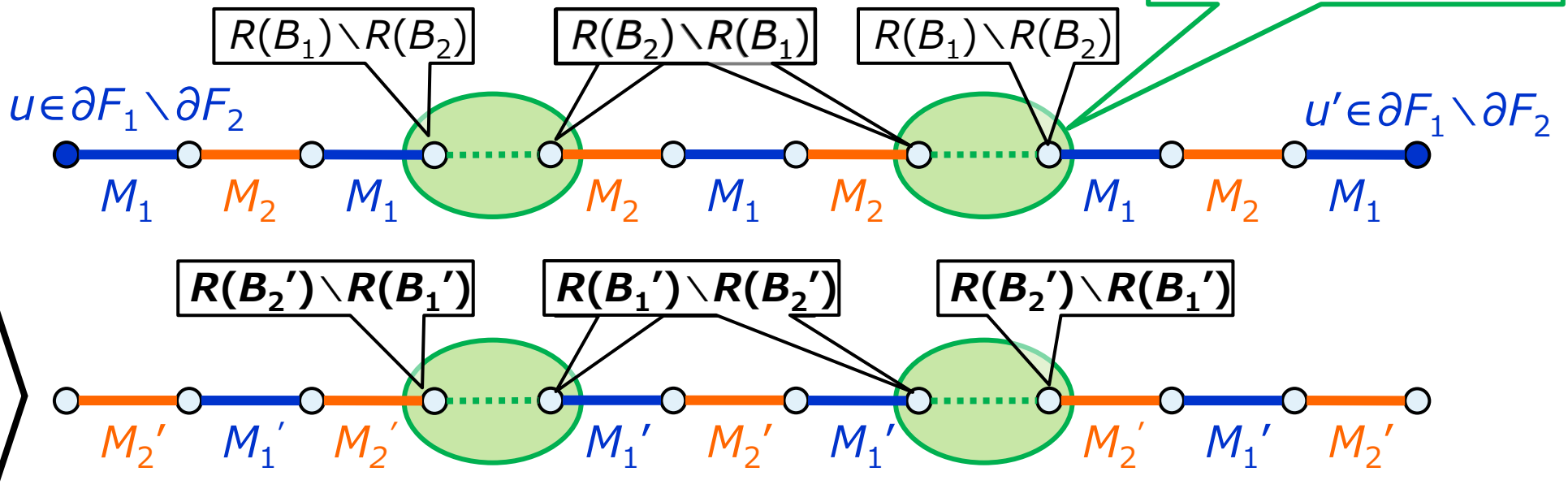
- Variant of **Edmonds' disjoint branchings theorem**
- All derives
 - TDI-ness of matching forests
 - Delta-matroid structure of matching forests
 - **Equitable partition into matching forests**
- Extended to **b -branchings**
 - Derives **equitable partition into b -branchings**



Proof Sketch for Matching Forests (1)

■ Suppose $|\partial F_1| - |\partial F_2| \geq 3 \rightarrow \exists u \in \partial F_1 \setminus \partial F_2$

■ **Case A**



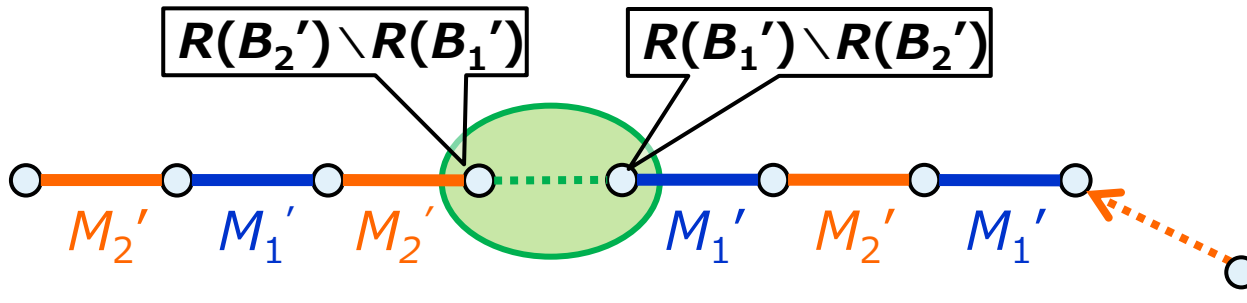
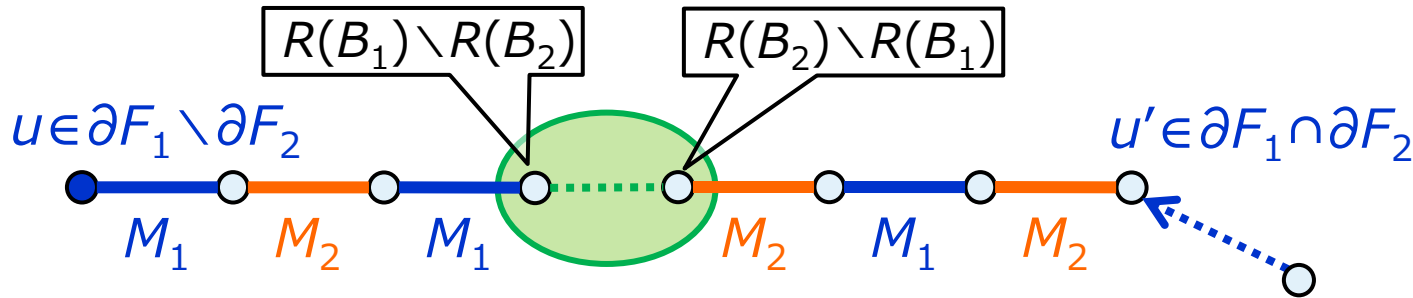
➤ Key Lemma shows the existence of B_1', B_2'

➤ $|\partial F_1| - |\partial F_2|$ decreases by **4**

Proof Sketch for Matching Forests (2)

■ Suppose $|\partial F_1| - |\partial F_2| \geq 3 \rightarrow \exists u \in \partial F_1 \setminus \partial F_2$

■ **Case B**

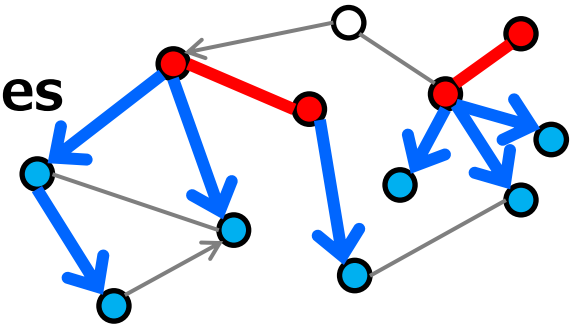


➤ Key Lemma shows the existence of B_1', B_2'

➤ $|\partial F_1| - |\partial F_2|$ decreases by **2**

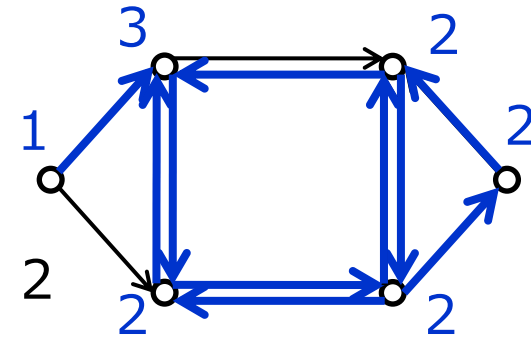
● Matching Forest

- Common generalization of matching and branching
- [Király, Yokoi '18] **Tricriteria equitability**
 - **Number of directed/undirected/all edges**
 - Cannot be optimized simultaneously
- [This talk] **Single-criterion equitability**
 - **Number of covered vertices**
 - Can always be optimized



● b -branching

- Generalization of branchings allowing indegree ≥ 2
- [This talk] **$(n+1)$ -criteria equitability**
 - **Number of all edges + indegree of each vertex v**
 - Can always be optimized simultaneously



END of Slides

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